

Frequency-smoothing encryption

Preventing snapshot attacks on deterministically encrypted data

Marie-Sarah Lacharité and Kenneth G. Paterson

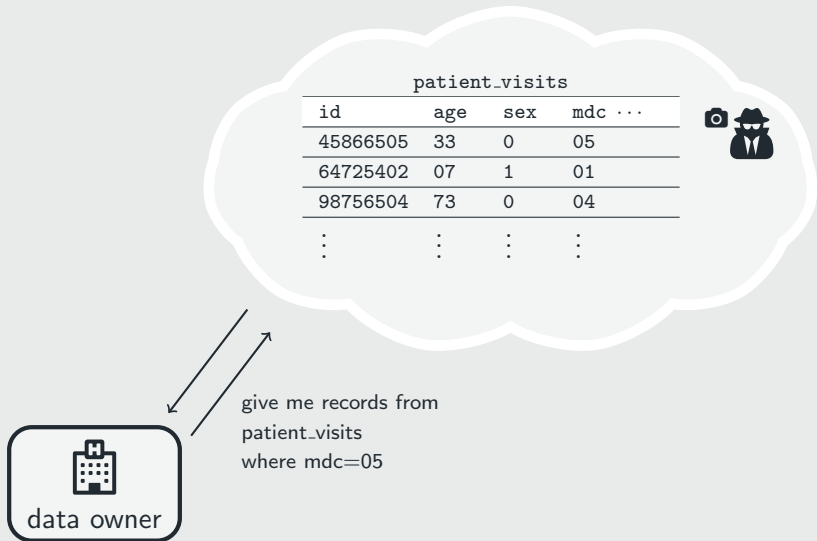
Information Security Group, Royal Holloway, University of London

6 March 2018

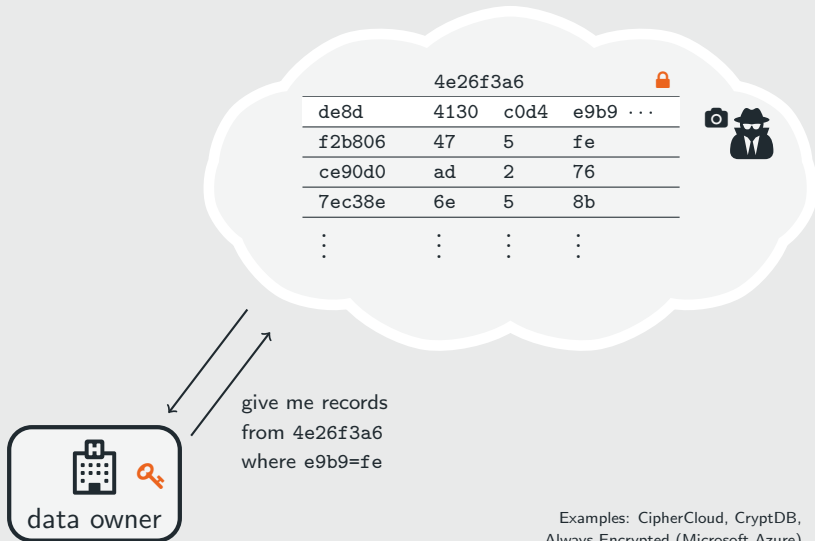
FSE 2018, Bruges



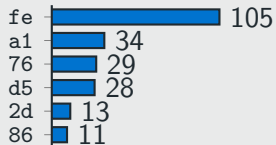
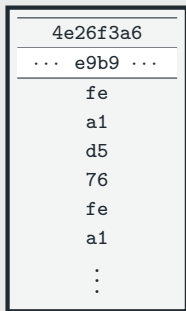
Outsourced database storage



Outsourced database storage with deterministic encryption



Inference attacks: an example

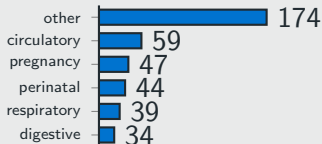
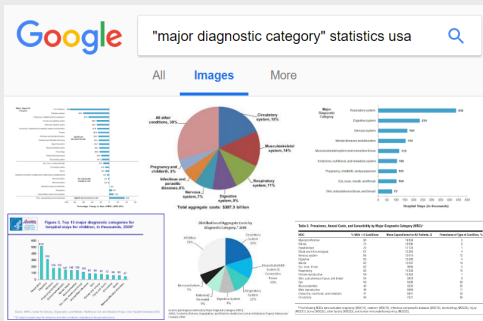
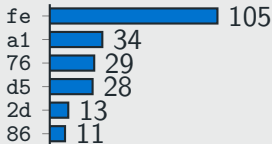


Inference attacks: an example

```

4e26f3a6
... e9b9 ...
-----
fe
a1
d5
76
fe
a1
:

```



[NKW15]

Inference Attacks on Property-Preserving Encrypted Databases

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recovered MDC values
in $\geq 20\%$ of records
for 75% hospitals

[GSB⁺17]

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Leakage-Abuse Attacks against Order-Revealing Encryption

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[PW16]

The Shadow Nemesis: Inference Attacks on Efficiently Deployable, Efficiently Searchable Encryption

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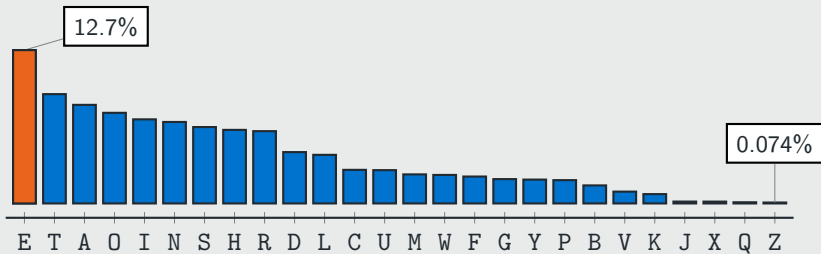
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Overview of our results

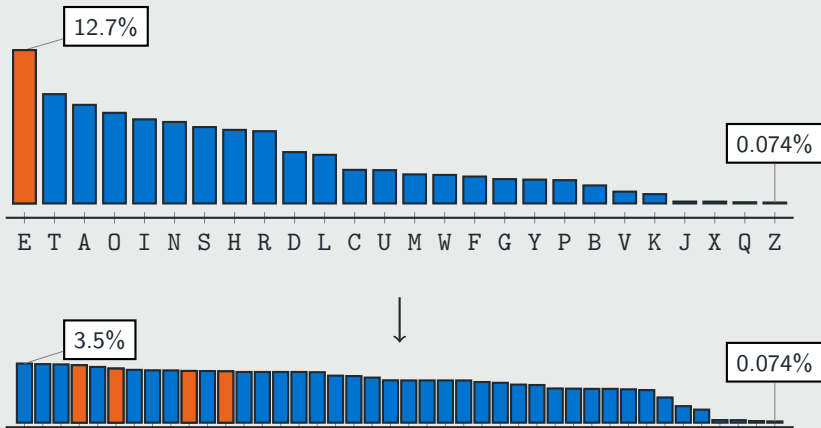
- frequency-smoothing (FS) encryption framework
- construction from homophonic encoding (HE) and deterministic encryption (DE)
- analytical and experimental evaluation of smoothness
- 8-bit FS encoding: recover $\geq 20\%$ of MDC values for only 2% of hospitals
 - when *exact* distribution is known

Frequency-smoothing encryption

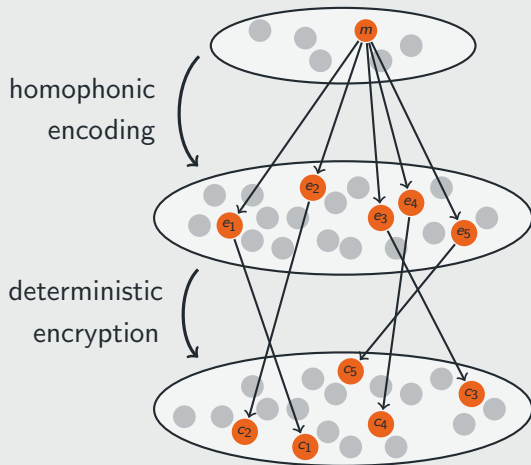
Inspiration: homophonic encoding (HE)



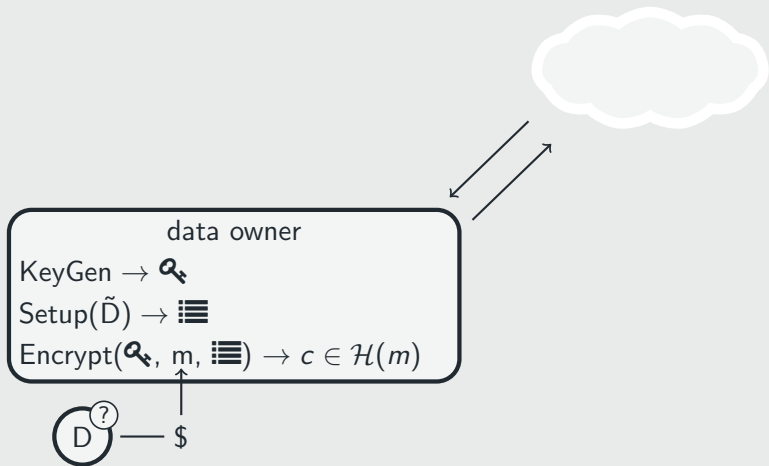
Inspiration: homophonic encoding (HE)



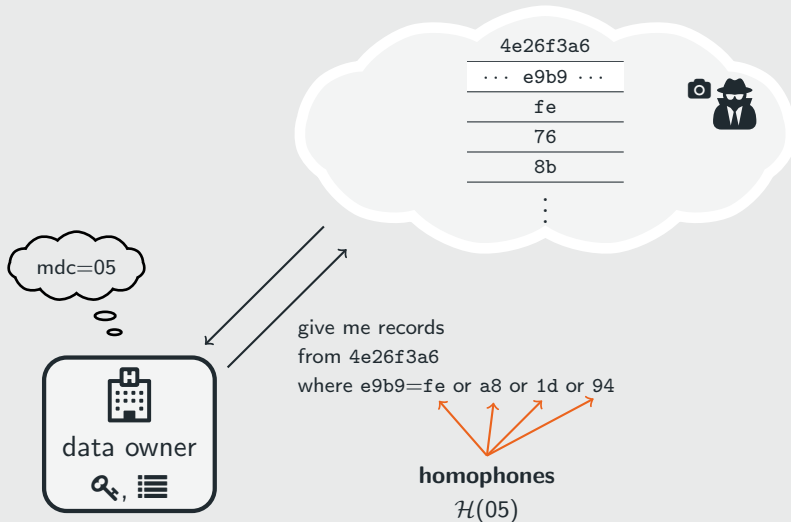
FS encryption from HE and DE



Frequency-smoothing (FS) encryption



Outsourced database storage with FS encryption



Frequency-smoothing (FS) encryption security

- adversary has its own estimate \hat{D} of the data's distribution
- *FS smoothness*: \mathcal{A} gets $\{c_1, \dots, c_N\}$, \tilde{D} , \hat{D}
 - are the N ciphertexts (i) real – generated by a FS encryption scheme with \tilde{D} , or (ii) fake – sampled from a set of size $|\mathcal{H}|$ uniformly at random?
- *FS message privacy*: \mathcal{A} gets $\{(m_1, c_1), \dots, (m_N, c_N)\}$, \tilde{D} , \hat{D}
 - are the N ciphertexts (i) real – generated by a FS encryption scheme with \tilde{D} , or (ii) fake – sampled from a set of size $|\mathcal{H}(m_i)|$ uniformly at random?

FS encryption from HE and DE: security

$$\left\{ \begin{array}{c} \text{HE smoothness} \\ + \\ \text{DE message privacy} \end{array} \right\} \Rightarrow \left\{ \begin{array}{c} \text{FS smoothness} \\ + \\ \text{FS message privacy} \end{array} \right\}$$

- *HE smoothness*: \mathcal{A} gets $\{e_1, \dots, e_N\}$, \tilde{D} , \hat{D}
 - are the N encodings (i) real – generated by an HE scheme with \tilde{D} , or (ii) fake – sampled from the set \mathcal{H} uniformly at random?
- *DE message privacy*: similar to IND\$ [Rog04]
 - could instantiate with small-domain PRP, format-preserving encryption, or synthetic IV mode [RS06]

HE smoothness when D is known

- distribution known by all: $D = \tilde{D} = \hat{D}$
 - so distribution D_e of encoded data depends only on D
- \mathcal{A} must distinguish D_e from uniform given N samples
- apply optimal distinguisher analysis from [BJV04]

Theorem

For any HE-SMOOTH adversary \mathcal{A} and sufficiently large N ,

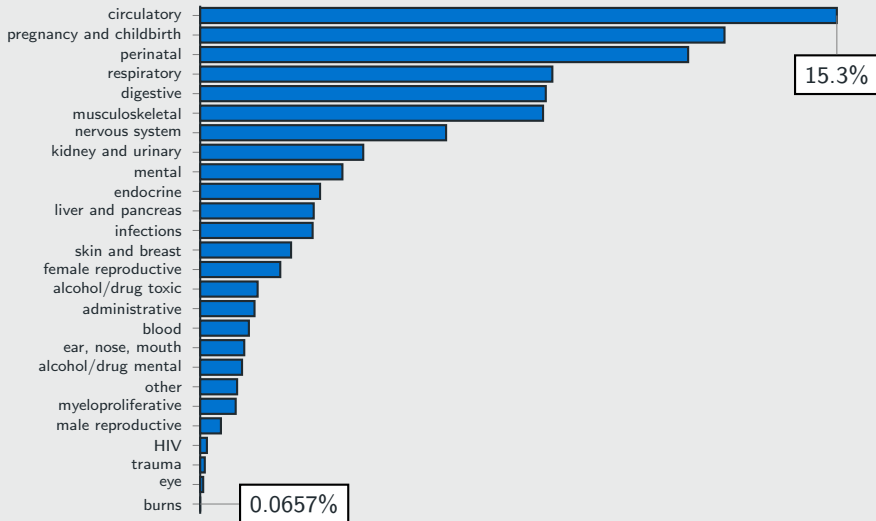
$$\text{Adv}(\mathcal{A}, D, N) \leq \left| \frac{1}{2} - \Phi \left(-\sqrt{\frac{N \cdot (\log |\mathcal{H}| - H_0(D_e))}{2}} \right) \right|$$

where $\Phi(\cdot)$ is cdf of the standard normal distribution and $H_0(\cdot)$ is Shannon entropy.

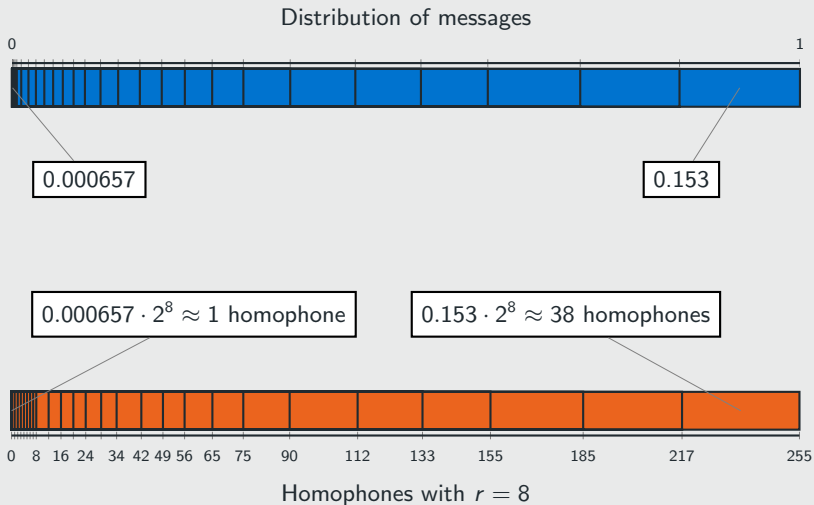
Interval-based homophonic encoding (IBHE)

- encodings are r -bit strings
- assign message m an interval of length $f_D(m) \cdot 2^r$
- choose homophones uniformly at random from this set
- maintain table of assigned intervals for decoding

IBHE example: MDC



IBHE example: MDC



IBHE example: MDC

- hospital has $N = 130\,000$ records
- probability of least frequent item is $2^{-11} \approx 0.00657$
- to limit smoothness advantage to $2^{-\epsilon}$, need encoding bitlength $r \approx 17 + \epsilon$
- main problem: query expansion

Practical security

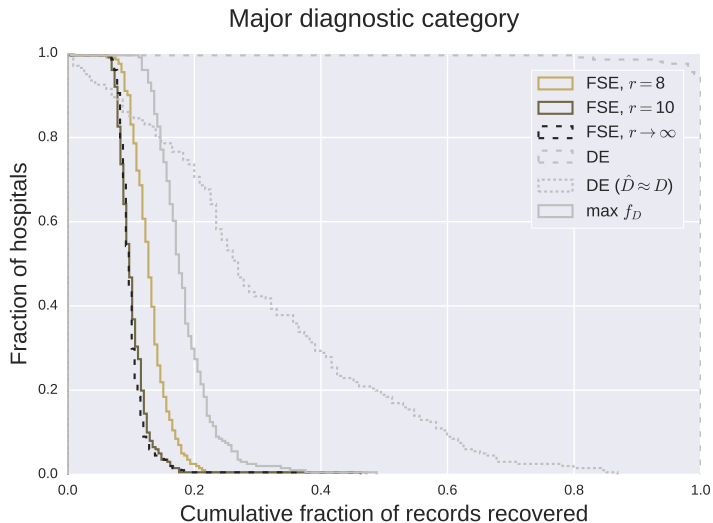
Experimental evaluation

- cryptographic security levels could require unacceptably large encoding lengths
 - and hence blow-up in query expansion
- empirically evaluate smoothness:
how many data items can adversary correctly decrypt?
- assume distribution D known by all
 - adversary knows how many homophones each message has
- what is optimal attack assuming only frequency information is meaningful?
 - message privacy easily achieved with a PRP

Maximum likelihood estimation (MLE)

- apply MLE to find most likely decryption function
- **MLE applied to deterministic encryption:** decrypt most frequent ciphertext to most frequent plaintext, and so on [LP15]
- **MLE applied to FS encryption:** decrypt $|\mathcal{H}(m_1)|$ most frequent ciphertexts to most frequent plaintext m_1 , and so on
- considers only “proper” decryption functions

Frequency-smoothing (FS) vs. deterministic (DE) encryption



Summary of contributions


- FS encryption thwarts snapshot inference attacks
- price to pay: query expansion, client storage
- see paper for
 - framework for *dynamic* FS schemes
 - FS construction from HE, PRF, and IV-based encryption
 - banded homophonic encoding scheme
- limited adversarial model, but part of **all** others

Summary of contributions


- FS encryption thwarts snapshot inference attacks
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