

# New techniques for trail bounds and application to differential trails in KECCAK

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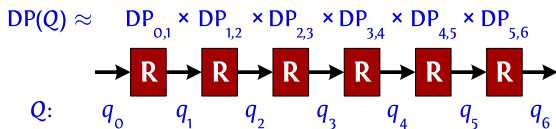
# Outline

- 1 Introduction
- 2 Generating trails
- 3 Scanning space of trails in  $\text{KECCAK-}f$
- 4 Experimental results
- 5 Conclusions

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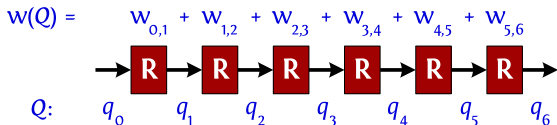
# Differential trails in iterated mappings



- ▶ Trail: the sequence of differences after each round
- ▶  $\text{DP}(Q)$ : fraction of pairs that exhibit  $q_i$  differences

# Differential trails and weight

$$w = -\log_2(DP)$$



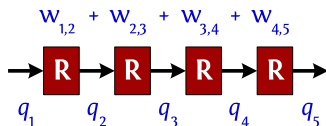
- ▶ The weight is the number of binary conditions that a pair must satisfy to exhibit  $q_i$  differences
- ▶ If *independent* conditions and  $w(Q) < b$ :  $\#\text{pairs}(Q) \approx 2^{b-w(Q)}$

# Trail extension

Given a trail, we can extend it

- ▶ forward: iterate over all differences  $R$ -compatible with  $q_5$
- ▶ backward: iterate over all differences  $R^{-1}$ -compatible with  $q_1$

Extension can be done recursively

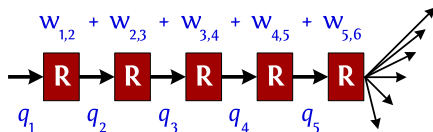


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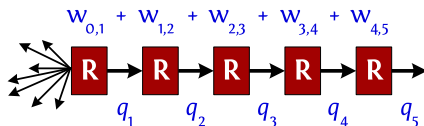


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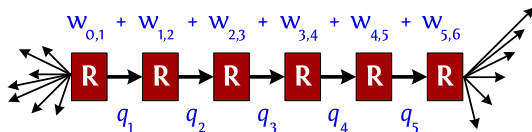


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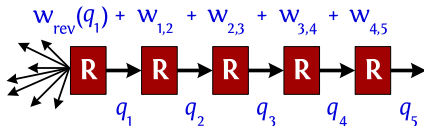


# Trail cores

- ▶ Minimum reverse weight:

$$w^{\text{rev}}(q_1) \triangleq \min_{q_0} w(q_0, q_1)$$

- ▶ Can be used to lower bound set of trails
- ▶ Trail core: set of trails with  $q_1, q_2, \dots$  in common



# Goals of this work

- ▶ Present general techniques to generate trails
- ▶ Improve bounds of differential trails in  $\text{KECCAK-}f$ 
  - ▶ By extending the space of trails in  $\text{KECCAK-}f$  that can be scanned with given computation resources

rounds	$\text{KECCAK-}f[200]$	$\text{KECCAK-}f[400]$	$\text{KECCAK-}f[800]$	$\text{KECCAK-}f[1600]$
2	8	8	8	8
3	20	this work	this work	32
4	46	this work	this work	this work
5	this work	this work	this work	this work
6	this work	this work	this work	this work

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# Generation of n-round trails of weight $\leq T$

## First-order approach

Starting from 1-round differentials with weight  $\leq \lfloor \frac{T}{n} \rfloor$

## Second-order approach

Starting from 2-round trails with weight  $\leq \lfloor \frac{2T}{n} \rfloor$

## Fact

The number of 2-round trails with weight  $\leq 2L$  is much smaller than the number of 1-round differentials with weight  $\leq L$ .

## Example: AES

AES has more than  $10^{11}$  round differentials with weight  $\leq 15$ , but no 2-round trail with weight  $\leq 30$

# Generating 2-round trails as tree traversal

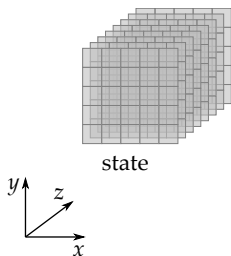
- ▶ 2-round trails are arranged in a tree
- ▶ Children are generated by adding groups of active bits without removing bits already added
- ▶ Pruning by lower bounding the weight of a node and its children

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# KECCAK- $f$

Operates on 3D state:



- ▶  $(5 \times 5)$ -bit **slices**
- ▶  $2^\ell$ -bit **lanes**
- ▶ parameter  $0 \leq \ell < 7$

Round function with 5 steps:

- ▶  $\theta$ : mixing layer
- ▶  $\rho$ : inter-slice bit transposition
- ▶  $\pi$ : intra-slice bit transposition
- ▶  $\chi$ : non-linear layer
- ▶  $\iota$ : round constants

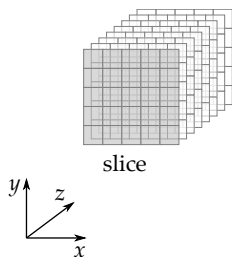
# rounds:  $12 + 2\ell$  for width  $b = 2^\ell 25$

- ▶ 12 rounds in KECCAK- $f[25]$
- ▶ 24 rounds in KECCAK- $f[1600]$



KECCAK- $f$ 

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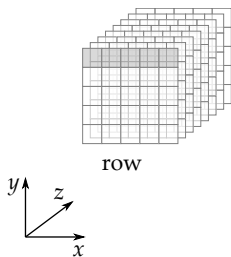
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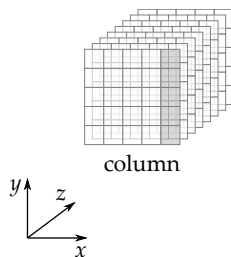
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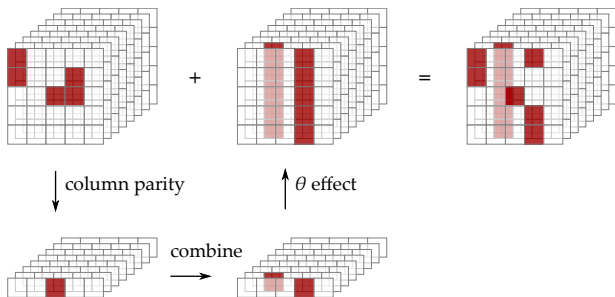
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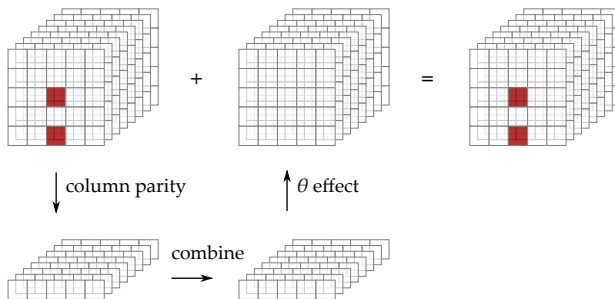
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# Properties of $\theta$



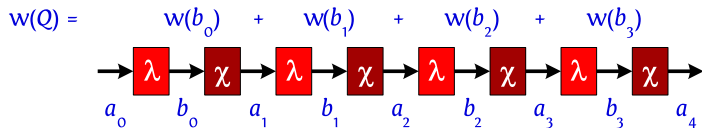
- ▶ The  $\theta$  map adds a pattern, that depends on the parity, to the state.
- ▶ **Affected** columns are complemented
- ▶ **Unaffected** columns are not changed

# The parity Kernel



- ▶  $\theta$  acts as the identity if parity is zero
- ▶ A state with parity zero is **in the kernel** (or in  $|K|$ )
- ▶ A state with parity non-zero is **outside the kernel** (or in  $|N|$ )

# Differential trails in KECCAK-f

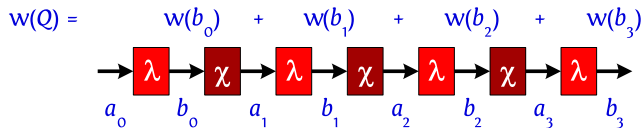


Round: linear step  $\lambda = \pi \circ \rho \circ \theta$  and non-linear step  $\chi$

- ▶  $a_i$  fully determines  $b_i = \lambda(a_i)$
- ▶  $\chi$  has degree 2:  $w(b_{i-1})$  independent of  $a_i$
- ▶ Minimum reverse weight:

$$w^{\text{rev}}(a_1) \triangleq \min_{b_0} w(b_0)$$

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# Differential trails in KECCAK-f

$$w(Q) = w_{\text{rev}}(a_1) + w(b_1) + w(b_2) + w(b_3)$$

The diagram illustrates a differential trail through a sequence of operations. It starts with an input  $a_1$  entering a red box labeled  $\lambda$ . The output of this box is  $b_1$ , which enters a dark red box labeled  $\chi$ . The output of  $\chi$  is  $a_2$ , which enters another red box labeled  $\lambda$ . The output of this second  $\lambda$  box is  $b_2$ , which enters a dark red box labeled  $\chi$ . The output of this second  $\chi$  box is  $a_3$ , which enters a final red box labeled  $\lambda$ . The output of this final  $\lambda$  box is  $b_3$ . Above the diagram, the weight of the trail is expressed as  $w(Q) = w_{\text{rev}}(a_1) + w(b_1) + w(b_2) + w(b_3)$ .

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# Covering the space of 3-round trail cores

$$w(Q) = w_{\text{rev}}(a_1) + w(b_1) + w(b_2)$$

$\rightarrow$   $\lambda$   $\rightarrow$   $\chi$   $\rightarrow$   $\lambda$   $\rightarrow$   
 $a_1$   $b_1$   $a_2$   $b_2$

- ▶ Space split based on parity of  $a_i$
- ▶ Four classes:  $|K|K|$ ,  $|K|N|$ ,  $|N|K|$  and  $|N|N|$

# Covering the space of 3-round trail cores

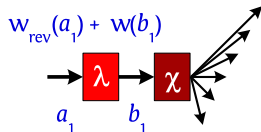
$$w(Q) = w_{\text{rev}}(a_1) + w(b_1)$$


The diagram shows a red square labeled with the Greek letter lambda (λ). An arrow points from the left into the square, and another arrow points from the square to the right. Below the left arrow is the label a<sub>1</sub>, and below the right arrow is the label b<sub>1</sub>.

- ▶ Generating  $(a_1, b_1)$
- ▶ Extending forward by one round

# Covering the space of 3-round trail cores

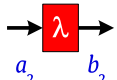
$w(Q) =$



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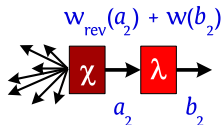
$$w_{\text{rev}}(a_2) + w(b_2)$$


The diagram shows a red square labeled with the Greek letter lambda ( $\lambda$ ). An arrow points from the left into the square, and another arrow points from the square to the right. Below the left arrow is the label  $a_2$ , and below the right arrow is the label  $b_2$ .

- ▶ Generating  $(a_2, b_2)$
- ▶ Extending backward by one round

# Covering the space of 3-round trail cores

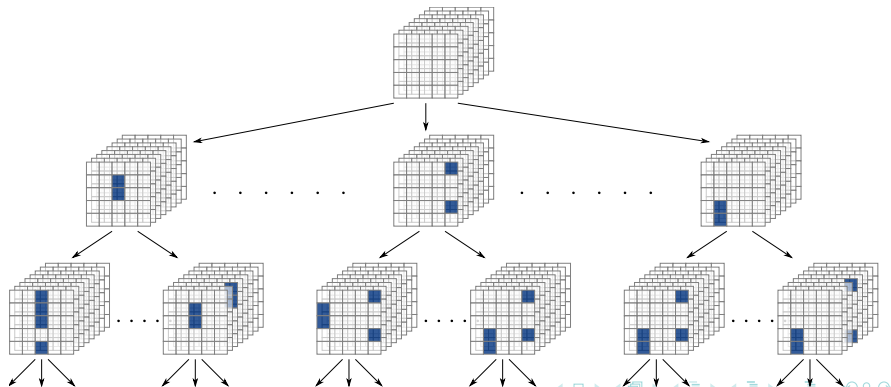
$w(Q) =$



- ▶ Generating  $(a_2, b_2)$
- ▶ Extending backward by one round

# Generating trail cores in $|K|$

- ▶ To stay in  $|K|$  units are *orbitals* = pairs of active bits in the same column
- ▶ A state  $a$  is a set of orbitals  $a = \{u_i\}_{i=1,\dots,n}$
- ▶ In the tree: the children of a node  $a$  are  $a \cup \{u_{n+1}\}$



# Order relation over units

- ▶ A total order relation over units allows avoiding duplicates
- ▶ With a total order  $\prec$  over units, a state is an ordered list of units:

$$a = (u_i)_{i=1,\dots,n} \text{ s.t. } u_1 \prec u_2 \prec \dots \prec u_n$$

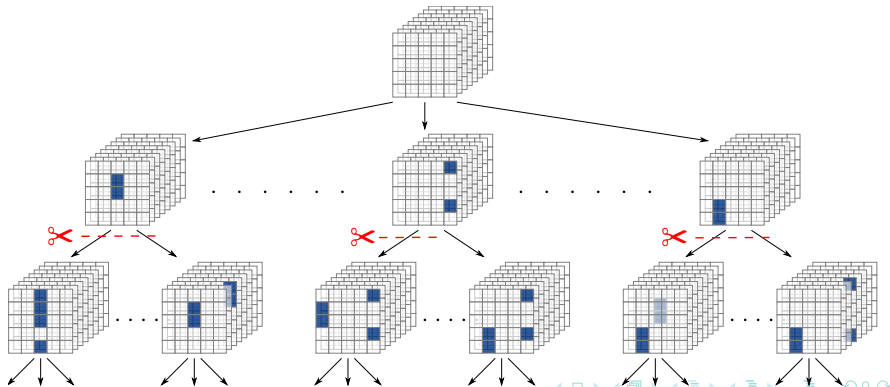
- ▶ In the tree: the children of a node  $a$  are

$$a \cup \{u_{n+1}\} \quad \forall u_{n+1} \text{ s.t. } u_n \prec u_{n+1}$$

- ▶ For orbitals: the lexicographic order  $[z, x, y_1, y_2]$

# Pruning by lower bounding the weight

- ▶ The weight is monotonic in the addition of orbitals
- ▶ The weight of  $a$  lower bounds the weight of all descendants of  $a$
- ▶ As soon as the search encounters  $a$  with weight above the limit,  $a$  and all its descendants can be safely pruned

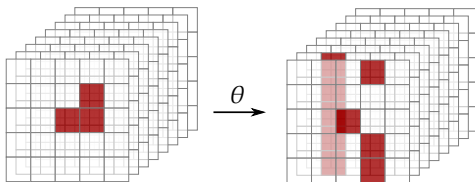




# Parity-bare states

Parity-bare state: a state with the minimum number of active bits before and after  $\theta$  for a given parity

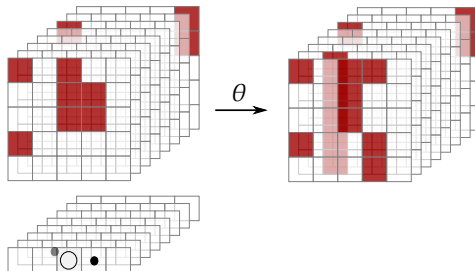
- ▶ 0 active bits in unaffected even columns
- ▶ 1 active bit in unaffected odd column
- ▶ 5 active bits in affected column either before or after  $\theta$



States in  $|N|$ 

## Lemma

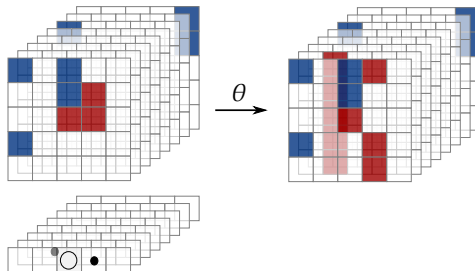
*Each state can be decomposed in a unique way in a parity-bare state and a list of orbitals*



States in  $|N|$ 

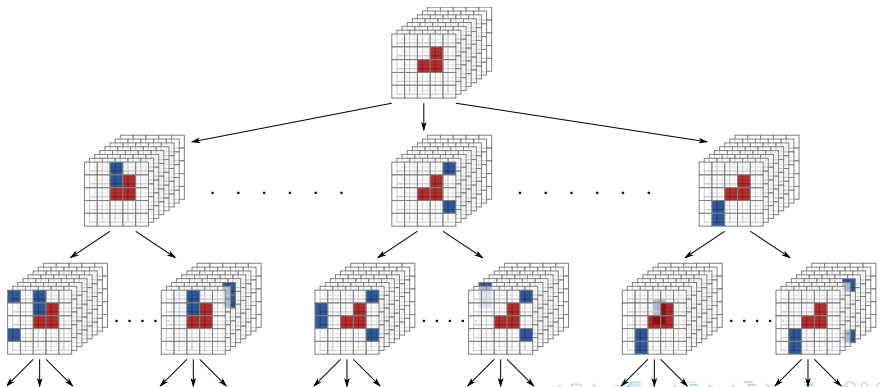
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*Each state can be decomposed in a unique way in a parity-bare state and a list of orbitals*



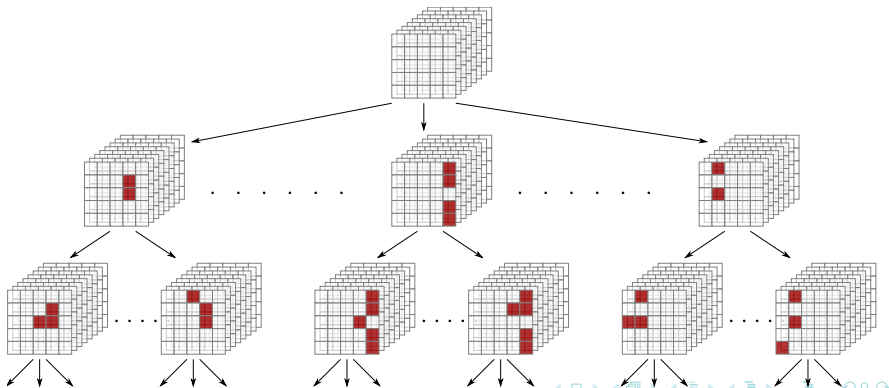
# Orbital tree

- ▶ Root: a parity-bare state
- ▶ Units: orbitals in unaffected columns
- ▶ Order: the lexicographic order on  $[z, x, y_1, y_2]$
- ▶ Bound: weight of the trail itself



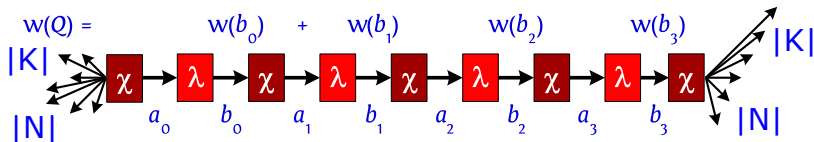
# Run tree

- ▶ Root: the empty state
- ▶ Units: column assignments
- ▶ Bound: by estimating maximum weight lost due to addition of new column assignments



# Trail extension

- ▶ forward: iterate  $a_4$  over all differences  $\chi$ -compatible with  $b_3$
- ▶ backward: iterate  $b_{-1}$  over all differences  $\chi^{-1}$ -compatible with  $a_0$
- ▶ in the kernel: restrict to differences with parity zero
- ▶ outside the kernel: restrict to differences with parity non-zero



# Forward extension

- ▶ Set of compatible states is an affine space  $\mathcal{A}(b_r) = e + V$
- ▶ Basis transformation:  $V = V_K + V_N$
- ▶ Extension in  $|K|$  by scanning  $e_K + V_K$ 
  - ▶ possible  $\Leftrightarrow e_K$  exists
- ▶ Extension in  $|N|$  by scanning  $e + V_K + V_N$
- ▶ Scanning as a tree traversal
  - ▶ root: is the offset
  - ▶ children: by incrementally adding basis vectors
  - ▶ bound: by estimating the maximum weight lost due to addition of basis vectors not already added

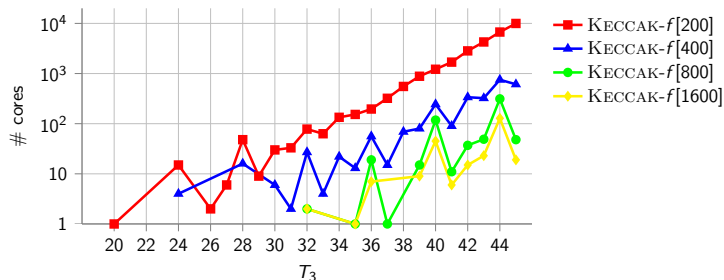
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# Experimental results

- ▶ All 3-round trail cores with weight  $\leq 45$



- ▶ No 6-round trail with weight  $\leq 91$

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# Conclusions

- ▶ General formalism to generate differential patterns as simple and efficient tree traversal
- ▶ New bounds for  $\text{KECCAK-}f$  and new trails with the lowest known weight

rounds	$b = 200$	$b = 400$	$b = 800$	$b = 1600$
2	8	8	8	8
3	20	24	32	32
4	46	[48,63]	[48,104]	[48,134]
5	[50,89]	[50,147]	[50,247]	[50,372]
6	[92,142]	[92,278]	[92,556]	[92,1112]

**Table:** Current bounds for the minimum weight of differential trails

Thanks for your attention